♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

♦ JohnsonBaugh's Algorithms, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm: Six cities

```
1 Foxville

O

2 Steger

O

3 Lusk
O
O
4 Springfield
O
5 Mystic
6 Del Rio
```

♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

1 Foxville 2 Steger

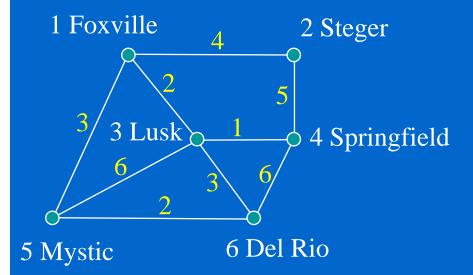
3 Lusk o o 4 Springfield

We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

5 Mystic6 Del Rio

♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

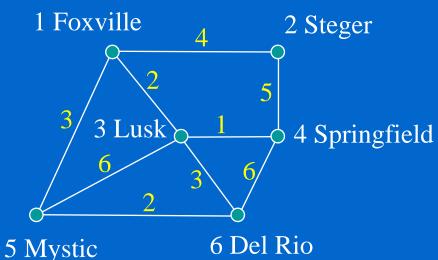
### Six cities



We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities



are minimized.

The estimated costs for some pairs of cities are as labeled.

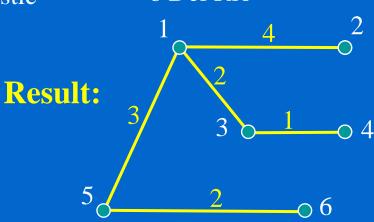
We want to construct a set of

interconnecting roads such

that one can reach any city

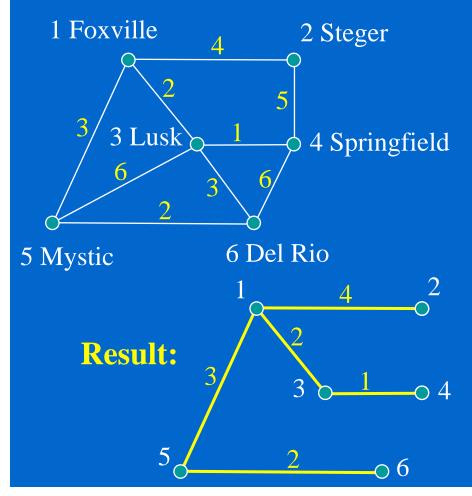
the total construction costs

from any starting city and

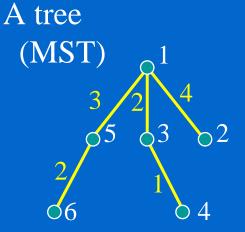


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

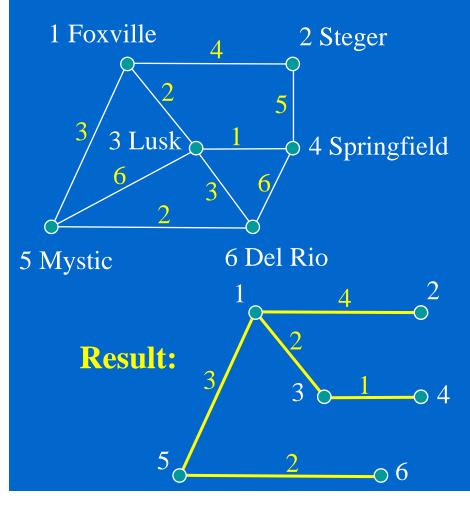


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

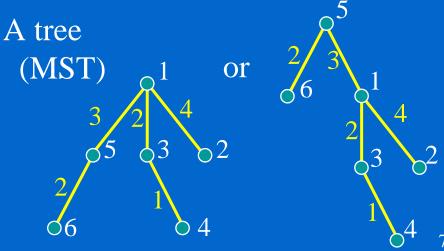


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

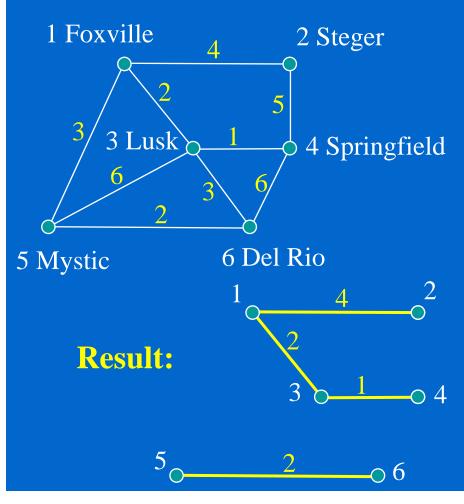


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

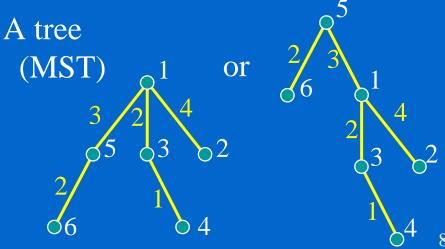


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

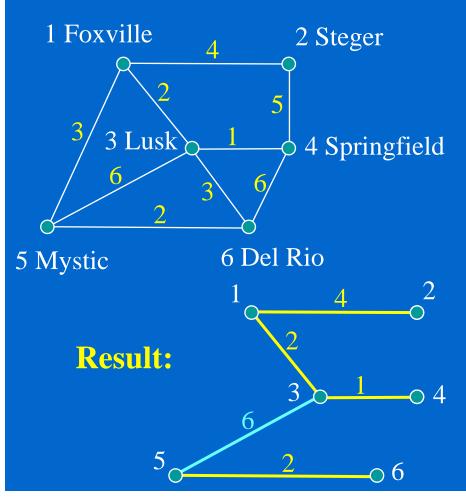


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

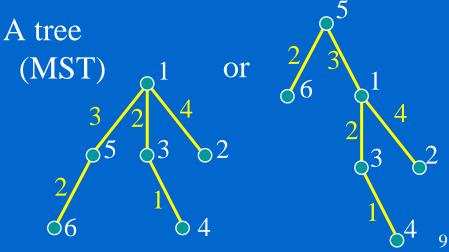


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

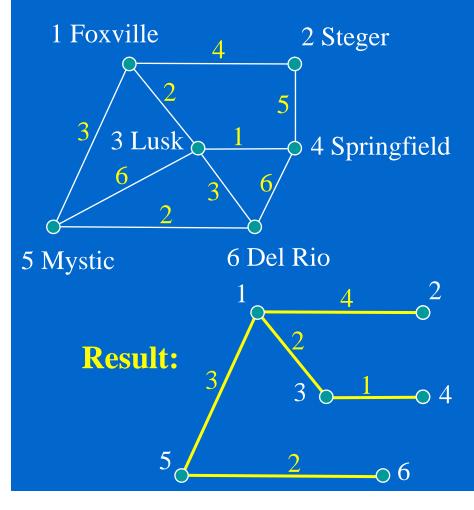


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

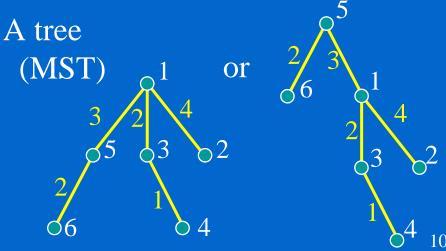


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

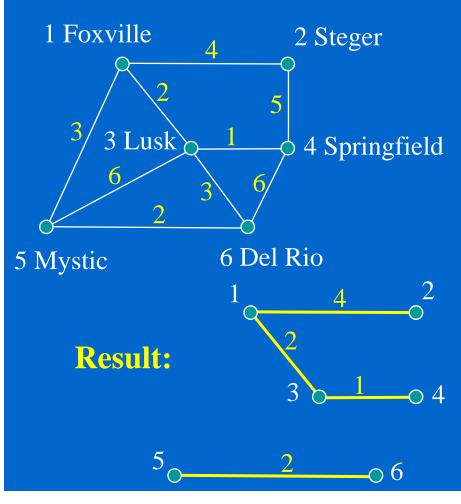


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

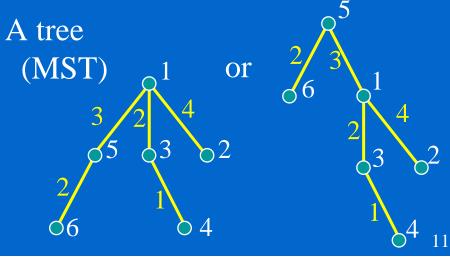


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

#### Six cities

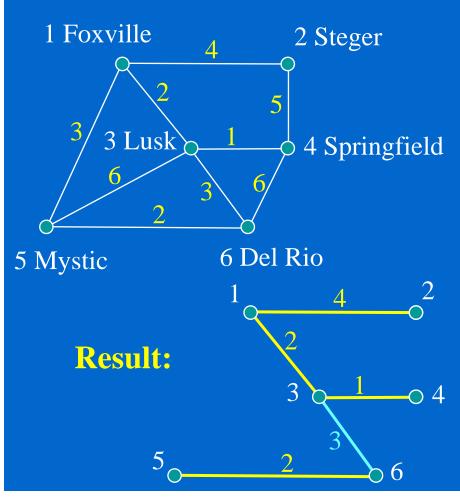


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

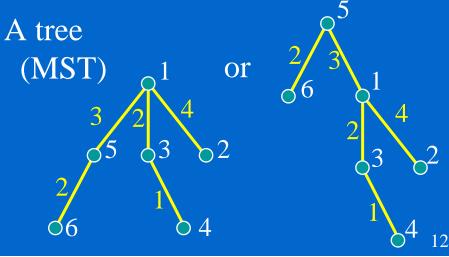


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

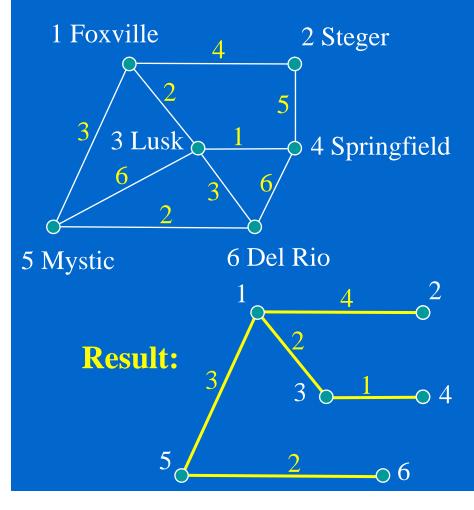


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

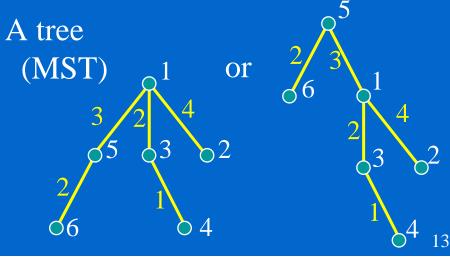


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

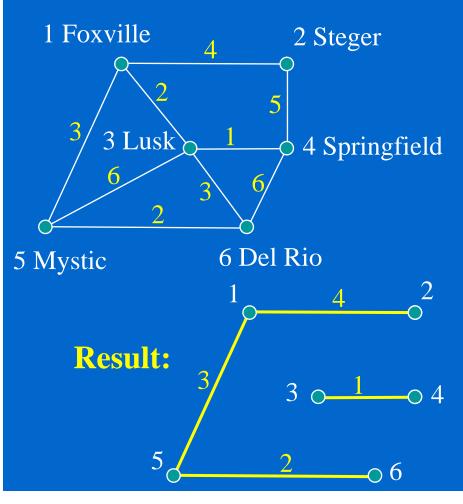


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

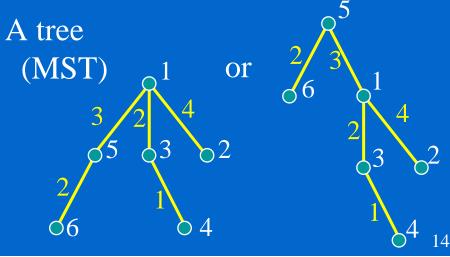


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

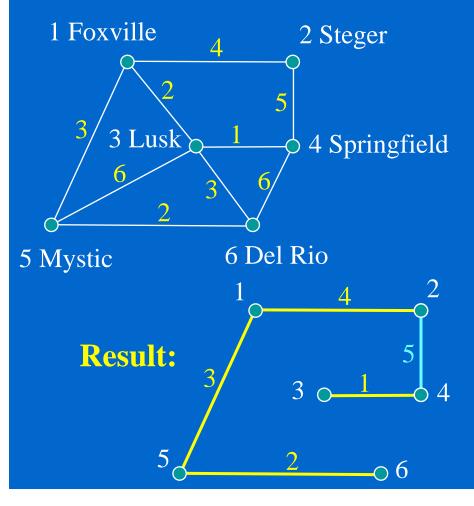


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

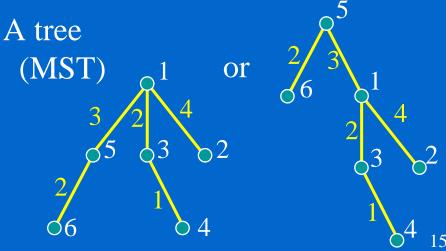


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

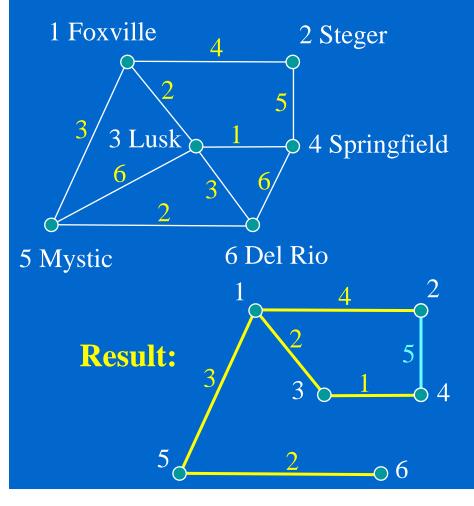


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

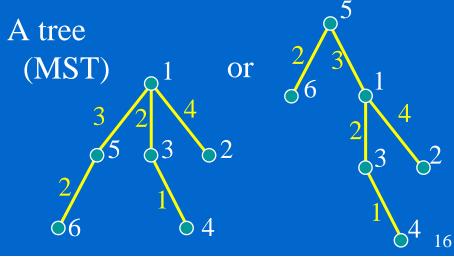


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

### Six cities

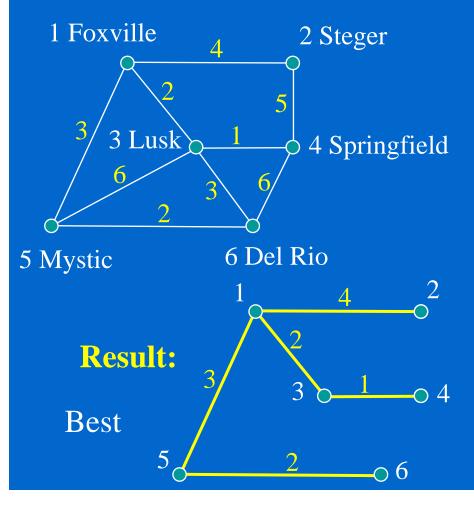


We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

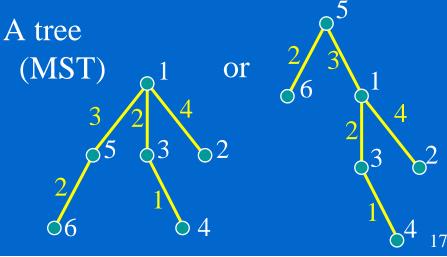


♦ JohnsonBaugh's *Algorithms*, Section 7.2 (page 275) find Minimal Spanning Tree (MST) with Kruskal's algorithm:

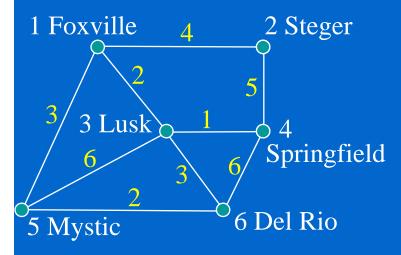
### Six cities

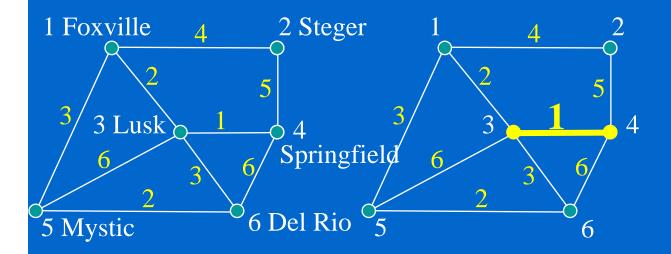


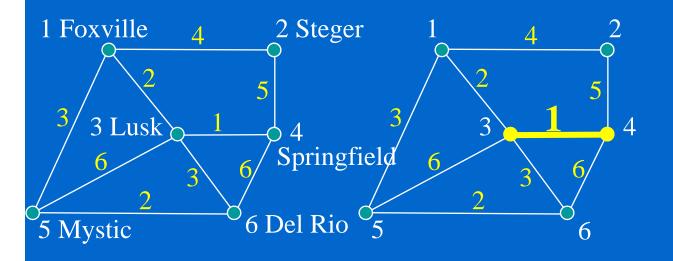
We want to construct a set of interconnecting roads such that one can reach any city from any starting city and the total construction costs are minimized.

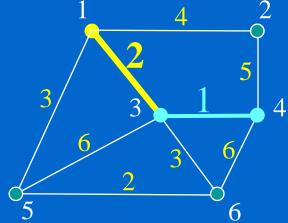


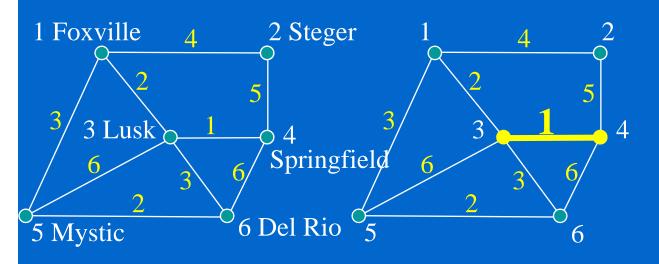


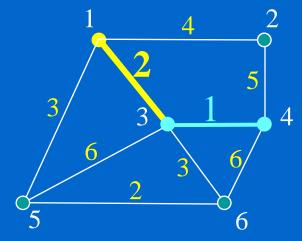


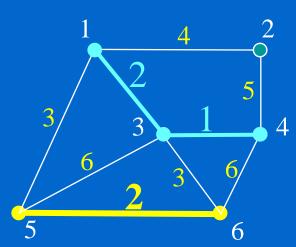


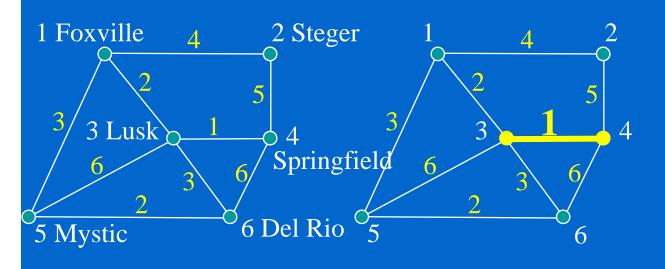


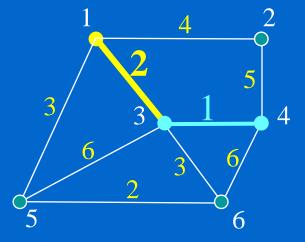


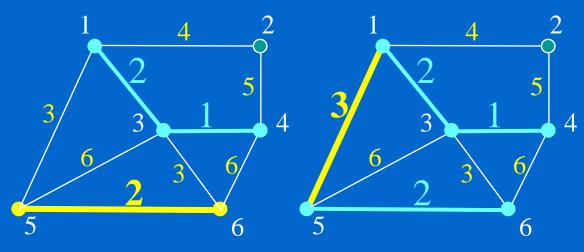


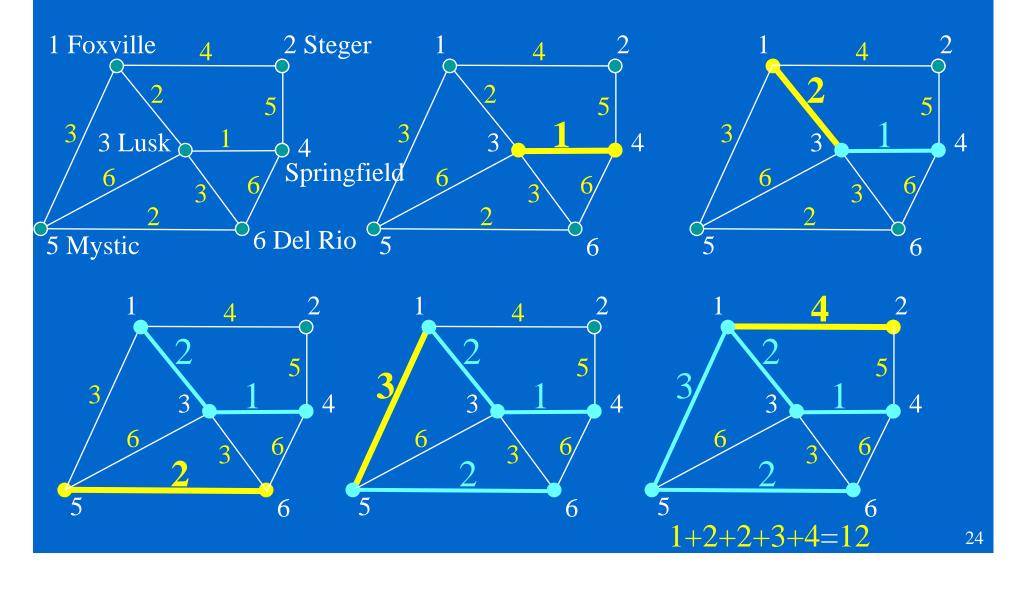


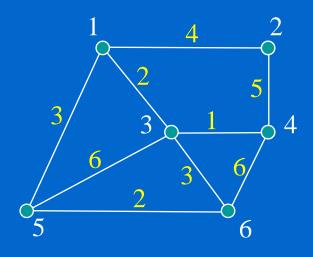




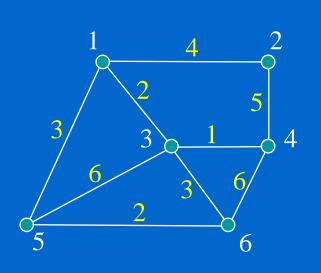








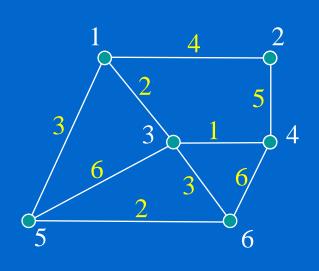
- find the edge with minimal weightadd to MST if the edge does not form a cycle



Array of edges:

(1,2,4),(1,3,2),(1,5,3),(2,4,5),(3,4,1),(3,5,6),(3,6,3),(4,6,6),(5,6,2)

- find the edge with minimal weight
- add to MST if the edge does not form a cycle



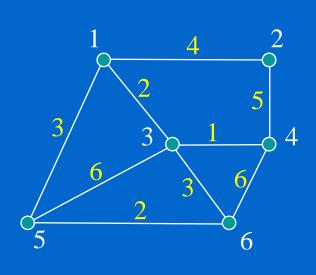
Array of edges:

$$(1,2,4),(1,3,2),(1,5,3),(2,4,5),(3,4,1),(3,5,6),$$
  
 $(3,6,3),(4,6,6),(5,6,2)$ 

Sorted array of edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 

- find the edge with minimal weight
- add to MST if the edge does not form a cycle



Array of edges:

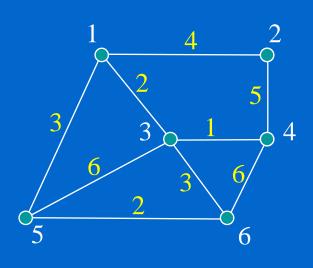
$$(1,2,4),(1,3,2),(1,5,3),(2,4,5),(3,4,1),(3,5,6),$$
  
 $(3,6,3),(4,6,6),(5,6,2)$ 

Sorted array of edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 

MST:{}

- 1 Ind the edge with minimal weight
- 2 add to MST if the edge does not form a cycle





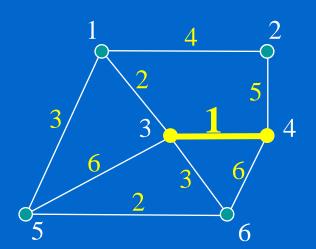
$$(1,2,4),(1,3,2),(1,5,3),(2,4,5),(3,4,1),(3,5,6),$$
  
 $(3,6,3),(4,6,6),(5,6,2)$ 

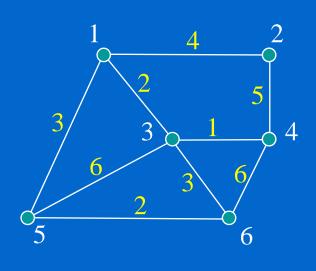
Sorted array of edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 

MST:{}

- 1 Ind the edge with minimal weight
- 2 add to MST if the edge does not form a cycle





Array of edges:

$$(1,2,4),(1,3,2),(1,5,3),(2,4,5),(3,4,1),(3,5,6),$$
  
 $(3,6,3),(4,6,6),(5,6,2)$ 

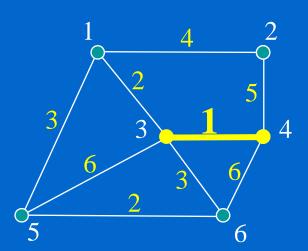
Sorted array of edges:

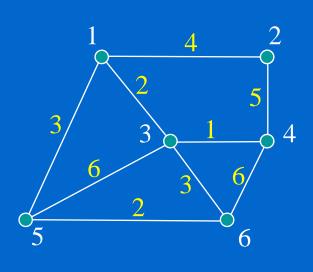
$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 

MST:{}

MST:{3,4}

- 1 Ind the edge with minimal weight
- 2 add to MST if the edge does not form a cycle





Array of edges:

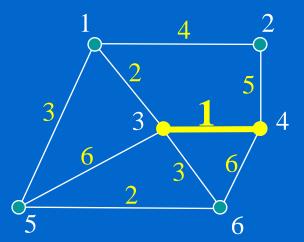
$$(1,2,4),(1,3,2),(1,5,3),(2,4,5),(3,4,1),(3,5,6),$$
  
 $(3,6,3),(4,6,6),(5,6,2)$ 

Sorted array of edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 

MST:{}

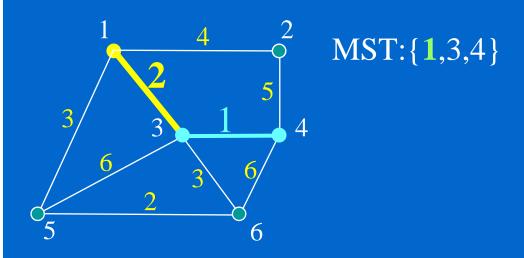
- find the edge with minimal weight add to MST if the edge does not form a cycle

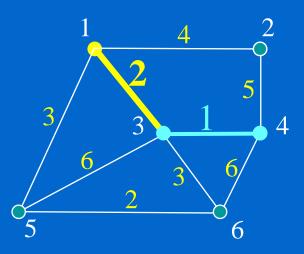


MST:{3,4}

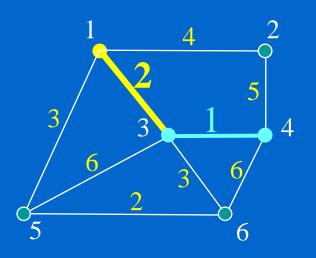
Remaining edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 





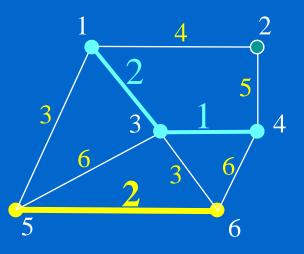
```
MST:{1,3,4}
Remaining edges:
(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),
(2,4,5),(3,5,6),(4,6,6)
```



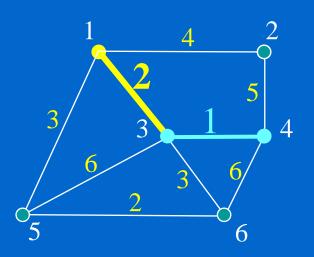
MST:{1,3,4}

Remaining edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$



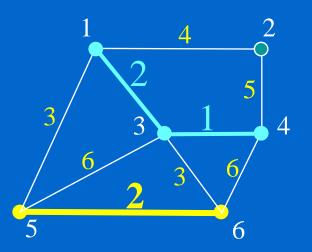
MST:{1,3,4},{5,6}



MST:{1,3,4}

Remaining edges:

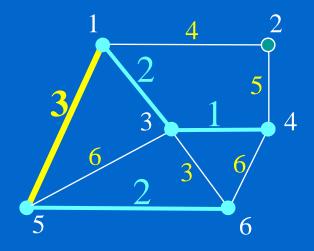
$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 



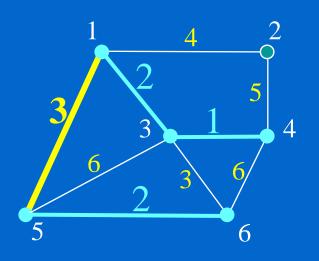
MST:{1,3,4},{5,6}

Remaining edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 



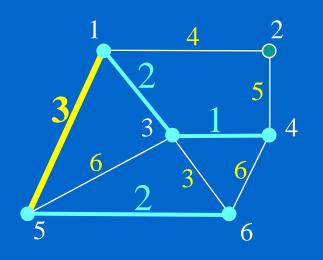
MST:{1,3,4,5,6}



```
MST:{1,3,4,5,6}
```

Remaining edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 

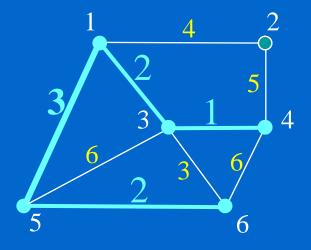


MST:{1,3,4,5,6}

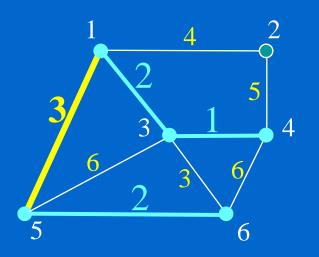
Remaining edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$

(2,4,5),(3,5,6),(4,6,6)



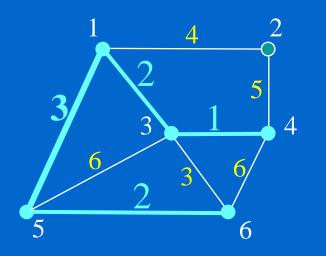
MST:{1,3,4,5,6}



MST:{1,3,4,5,6}

Remaining edges:

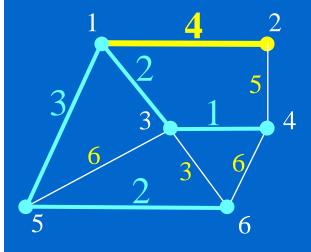
$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 



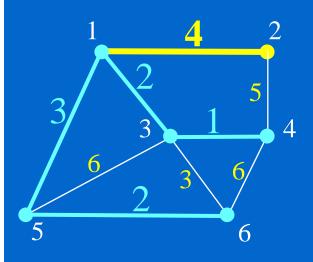
MST:{1,3,4,5,6}

Remaining edges:

$$(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),$$
  
 $(2,4,5),(3,5,6),(4,6,6)$ 



MST:{1,2,3,4,5,6}



MST:{1,2,3,4,5,6}

Remaining edges:

(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),

(2,4,5),(3,5,6),(4,6,6)

```
Array of edges: (vertex1, vertex2, weight)
    (1,2,4),(1,3,2),(1,5,3),(2,4,5),(3,4,1),(3,5,6),(3,6,3),(4,6,6),(5,6,2)
♦ Implementation ①: 2-dimensional arrays (or parallel arrays)
       int edges[][3] = \{\{1,2,4\},\{1,3,2\},\{1,5,3\},\{2,4,5\},\{3,4,1\},
                         {3,5,6},{3,6,3},{4,6,6},{5,6,2};
       int nEdges = sizeof(edges) / sizeof(int[3]);
  Implementation 2: array of struct
       struct Edge {
         int vertex1, vertex2, weight;
       struct Edge edges[] = \{\{1,2,4\},\{1,3,2\},\{1,5,3\},\{2,4,5\},\{3,4,1\},
                               {3,5,6},{3,6,3},{4,6,6},{5,6,2};
       int nEdges = sizeof(edges) / sizeof(struct Edge);
                                                                           42
```

### **Sorted array of edges:**

```
(3,4,1),(5,6,2),(1,3,2),(1,5,3),(3,6,3),(1,2,4),(2,4,5),(3,5,6),(4,6,6)
```

```
    Simple selection sort on
    2-dimensional arrays
    (slightly different results
    from previous slides)
```

```
void swap(int a[3], int b[3]) {
  int tmp, i;
  for (i=0; i<3; i++) {
    tmp = a[i];
    a[i] = b[i];
    b[i] = tmp;
  }
}</pre>
```

```
01 void selectionSort(int edges[][3], int nEdges) {
     int i, max;
02
03
     for (i=0; i<nEdges-1; i++) {
04
        max = findMaximum(edges, nEdges-i);
05
        swap(edges[nEdges-i-1], edges[max]);
06
07 }
08
09 int findMaximum(int edges[][3], int nEdges) {
     int i, max=nEdges-1;
10
     for (i=nEdges-2; i>=0; i--)
       if (edges[i][2] > edges[max][2])
12
          max = i;
13
14
     return max;
15 }
```

### **Sorted array of edges:**

```
(3,4,1),(5,6,2),(1,3,2),(1,5,3),(3,6,3),(1,2,4),(2,4,5),(3,5,6),(4,6,6)
```

stdlib qsort on array of structs

```
#include <stdlib.h>
int compare(void *arg1, void *arg2) {
   return ((struct Edge *)arg1)->weight - ((struct Edge *)arg2)->weight;
}
qsort(edgelist, nEdges, sizeof(struct Edge), compare);
```

### **Sorted array of edges:**

```
(3,4,1),(1,3,2),(5,6,2),(1,5,3),(3,6,3),(1,2,4),(2,4,5),(3,5,6),(4,6,6)
```

→ requires a stable sorting algorithm: e.g. bubble, bucket, insertion, counting, merge, radix, ...

 $MST:\{\} \rightarrow \{3,4\} \rightarrow \{1,3,4\} \rightarrow \{1,3,4\}, \{5,6\} \rightarrow \{1,3,4,5,6\} \rightarrow \{1,2,3,4,5,6\}$ 

- ♦ Require "set processing" tools: union, comparison
- ♦ Specially, these are disjoint sets (Section 3.6 of JohnsonBaugh, pp.150):
  - \* Set members are held in the same tree, root node represents the set
  - \* use an array *parent* to implement the 1 3 set membership and provide three interfaces:
    - **\* makeset**(i): construct the set {i}
    - **♦ findset**(i): returns the representative node of the set
    - ★ union(i,j): joins the set containing i and the set containing j

```
void makeset(int i, int nNodes, int parent[]) {
  if ((i<0)||(i>=nNodes)) return;
  parent[i] = i;
}
```

```
int findset(int i, int nNodes, int parent[]) {
   if ((i<0)||(i>=nNodes)) return -1;
   while (i != parent[i])
      i = parent[i];
   return i;
}
```

 $MST:\{\} \rightarrow \{3,4\} \rightarrow \{1,3,4\} \rightarrow \{1,3,4\}, \{5,6\} \rightarrow \{1,3,4,5,6\} \rightarrow \{1,2,3,4,5,6\}$ 

- ♦ Require "set processing" tools: union, comparison
- ♦ Specially, these are disjoint sets (Section 3.6 of JohnsonBaugh, pp.150):
  - \* Set members are held in the same tree, root node represents the set
  - \* use an array *parent* to implement the 1 3 4 set membership and provide three interfaces:
    - **\* makeset**(i): construct the set {i}
    - **♦ findset**(i): returns the representative node of the set
    - ★ union(i,j): joins the set containing i and the set containing j

parent

```
    1
    2
    3
    4
    5
    6

    5
    1
    5
    5
```

```
\label{eq:continuous_parent} \begin{tabular}{ll} void $\mathsf{makeset}(int i, int nNodes, int parent[]) $\{$ if $((i<0)||(i>=nNodes))$ return; \\ parent[i] = i; \\ \end{tabular} int $\mathsf{findset}(int i, int nNodes, int parent[]) $\{$ if $((i<0)||(i>=nNodes))$ return $-1$; \\ while $(i != parent[i])$ \\ $i = parent[i]$; \\ $return $i$; \\ \end{tabular}
```

```
void mergetrees(int i, int j, int nNodes, int parent[]) {
  if (((i < 0) || (i > = nNodes)) || ((j < 0) || (j > = nNodes))) return;
  parent[i] = j;
void union(int i, int j, int nNodes, int parent[]) {
  if (((i<0)||(i>=nNodes))||((j<0)||(j>=nNodes))) return;
  mergetrees(findset(i, nNodes, parent), findset(j, nNodes, parent), nNodes, parent);
                                                    • find the edge with minimal weight
                                                    2 add to MST if the edge does not
                                                       form a cycle
for (iEdge=0,treeSize=0; treeSize<nNodes; iEdge++) {
     if (findset(edgelist[iEdge][0], nNodes, nodeSet) !=
         findset(edgelist[iEdge][1], nNodes, nodeSet)) {
        totalWeight = totalWeight + edgelist[iEdge][2]; treeSize++;
        union(edgelist[iEdge][0], edgelist[iEdge][1], nNodes, nodeSet);
```