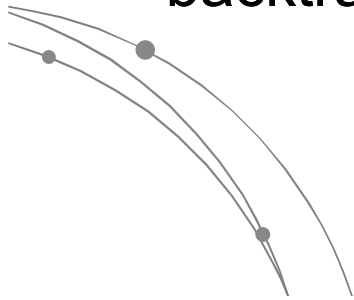


Permutation, Combination, and Related Problems

- backtracking algorithms



Pei-yih Ting

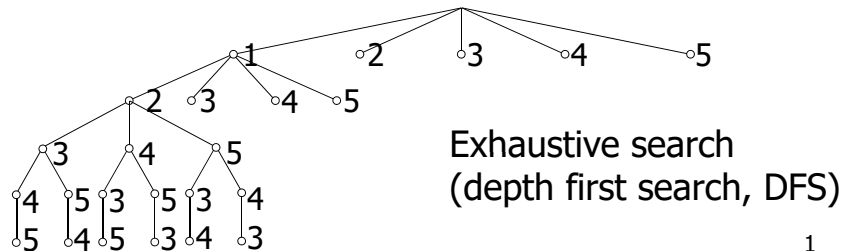
1

Introduction

- You might need to generate permutations, combinations, partitions to enumerate all possible configurations in order to find the optimal solutions or brute-force solutions of some problems
- Procedural programming is about data processing
- To design the program:
 - figure out how the data/configuration changes
 - find the suitable representation of data in a program
 - figure out the “process” that transforms the data step by step

2

Generating Permutations



5! = 120
permutations

1 2 3 4 5
1 2 3 5 4
1 2 4 3 5
1 2 4 5 3
1 2 5 3 4
1 2 5 4 3
...

				1
			1	2
		1	2	3
1	2	3	4	5
2	3	4	5	
3	4	5		
4	5			
5				

3

n-layer for loop Implementation

```
int i1, i2, i3, i4;
for (i1=1; i1<=4; i1++) {
    for (i2=1; i2<=4; i2++) {
        if (i2 == i1) continue;
        for (i3=1; i3<=4; i3++) {
            if ((i3==i2)|| (i3==i1)) continue;
            for (i4=1; i4<=4; i4++) {
                if ((i4==i3)|| (i4==i2)|| (i4==i1)) continue;
                printf("%d %d %d %d\n", i1, i2, i3, i4);
            }
        }
    }
}
```

i1	i2	i3	i4
----	----	----	----

This is a quick implementation but **NOT scalable**.

4

2-layer for loop Implementation

- Consider this simplified loop without collision constraints

```
int data[4];
for (data[0]=1; data[0]<=4; data[0]++) {
    for (data[1]=1; data[1]<=4; data[1]++) {
        for (data[2]=1; data[2]<=4; data[2]++) {
            for (data[3]=1; data[3]<=4; data[3]++) {
                printf("%d %d %d %d\n", data[0],
                    data[1], data[2], data[3]);
            }
        }
    }
}
```

data	i1	i2	i3	i4
	1	1	1	1
	1	1	1	2
	1	1	1	3
	1	1	1	4
	1	1	2	1
	1	1	2	2
	1	1	2	3
	1	1	2	4
	1	1	3	1

- It is still not scalable. ...
- Is there a **scalable** program structure that generates the same (i1, i2, i3, i4) counting-up sequence? yes

5

Equivalent 2-layer for loop

```
void nextIndex(int index[], int n) {
    int i;
    for (i=n-1; i>0; i--)
        if (index[i]<n)
            { index[i]++; return; }
    else
        index[i] = 1;
    index[0]++;
}
```

index	1	1	1	1
	1	1	1	2
	1	1	1	3
	1	1	1	4
	1	1	2	1
	1	1	2	2
	1	1	2	3
	1	1	2	4
	1	1	3	1

```
int i, index[4], n=4;
for (i=0; i<n; i++)
    index[i] = 1;
for (; index[0]<=n; nextIndex(index, n))
    printArray(index, n);
```

Scalable

6

2-layer for loop for Permutation

```
int index[4], n=4;
for (i=0; i<n; i++)
    index[i] = i+1;
for (; index[0]<=n; nextIndex(index, n))
    if (isValid(index, n))
        printPerm(index, n);
```

index	1	2	3	4
	1	2	4	3
	1	3	2	4
	1	3	4	2
	1	4	2	3
	1	4	3	2
	2	1	3	4
	2	1	4	3

```
int isValid(int index[], int n) {
    int i, j;
    for (i=0; i<n-1; i++)
        for (j=i+1; j<n; j++)
            if (index[i]==index[j])
                return 0;
    return 1;
}
```

```
void nextIndex(int index[], int n) {
    int i;
    for (i=n-1; i>0; i--)
        if (index[i]<n)
            { index[i]++; return; }
    else
        index[i] = 1;
    index[0]++;
}
```

7

```
int next(int perm[], int used[], int n) {
    int i, j, k;
    for (i=n-1; i>=0; i--) {
        used[perm[i]] = 0;
        do
            perm[i]++;
        while ((perm[i]<=n)&&used[perm[i]]);
        if (perm[i]<=n) {
            used[perm[i]] = 1;
            for (k=0; k<perm[i]; k++) {
                while (used[k]);
                perm[k] = k, used[k] = 1;
            }
            return 1;
        }
    }
    return 0;
}
```

Counting w/o invalid() check

perm	1	2	3
	1	3	2
	2	1	3
	2	3	1
	3	1	2
	3	2	1

```
int perm[10], used[11], n=3;
for (i=0; i<n; i++)
    perm[i] = i+1, used[i+1] = 1;
do
    printArray(perm, n);
while (next(perm, used, n));
```

8

Another Implementation

```
void print(int m, int num[]);
void next(int n, int num[], int it);
int main() {
```

```
    int i, n=4, num0[] = {0, 0, 1, 2, 3}, *num=num0+1;
```

```
    while (num[-1]==0) {
```

```
        print(n, num);
```

```
        i=n-1;
```

```
        do
```

```
            next(n, num, --i);
```

```
        while (i>=0&&num[i]==n);
```

```
        while (++i<n)
```

```
            num[i]=-1, next(n, num, i);
```

```
    } return 0;
```

```
}
```

```
void print(int n, int num[]) {
    for (int i=0; i<n; i++)
        printf("%d ", num[i]);
    printf("\n");
}
```

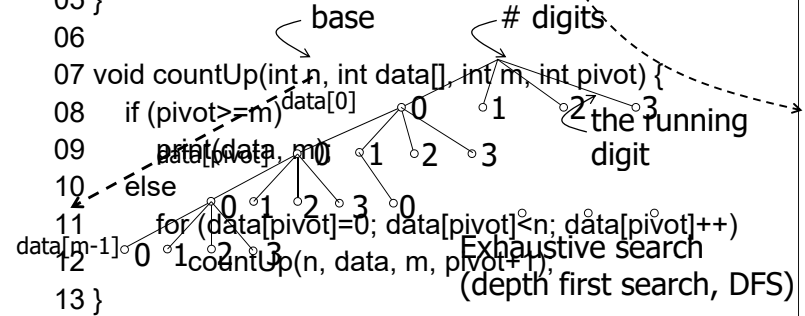
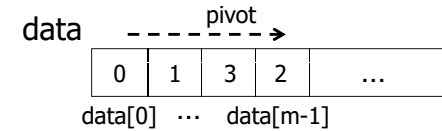
```
void next(int n, int num[], int it) {
    int i;
    for (num[it]++; num[it]<n; num[it]++) {
        for (i=0; i<it; i++)
            if (num[i]==num[it]) break;
        if (i==it) return;
    }
}
```

找到比 **num[it]** 大的下一個可用數字, 找不到 **num[it]** 就設為 **n**

		num[0]			
num[-1]					num[n-1]
0	0	1	2	3	
	0	1	3	2	
	0	2	1	3	
	0	2	3	1	
	0	3	1	2	
	0	3	2	1	
	1	0	2	3	
			...		

Recursive Counting Up

```
01 int main() {
02     int data[10];
03     countUp(4, data, 4, 0);
04     return 0;
05 }
06
07 void countUp(int n, int data[], int m, int pivot) {
08     if (pivot>=m)
09         print(data, m);
10     else
11         for (data[pivot]=0; data[pivot]<n; data[pivot]++)
12             countUp(n, data, m, pivot+1);
13 }
```

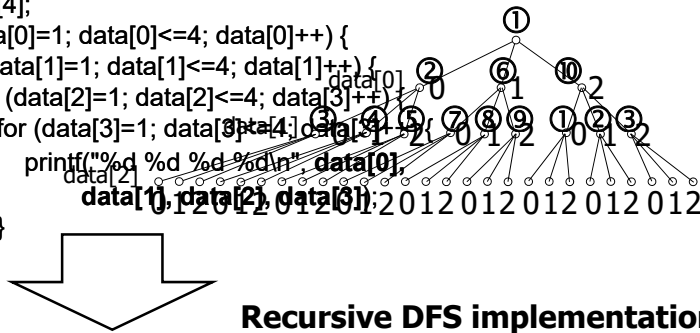


```
0 0 0 0
0 0 0 1
0 0 0 2
0 0 0 3
0 0 1 0
...
0 0 3 3
0 1 0 0
0 1 0 1
0 1 0 2
0 1 0 3
0 1 1 0
...
0 1 3 2
0 1 3 3
0 2 0 0
...
1 0 0 0
...
3 3 3 2
3 3 3 3
```

10

Generalization of Deep for Loops

```
01 int data[4];
02 for (data[0]=1; data[0]<=4; data[0]++) {
03     for (data[1]=1; data[1]<=4; data[1]++) {
04         for (data[2]=1; data[2]<=4; data[2]++) {
05             for (data[3]=1; data[3]<=4; data[3]++) {
06                 print("%d %d %d %d\n", data[0],
07                     data[1], data[2], data[3]);
08             }
09         }
10     }
11 }
```



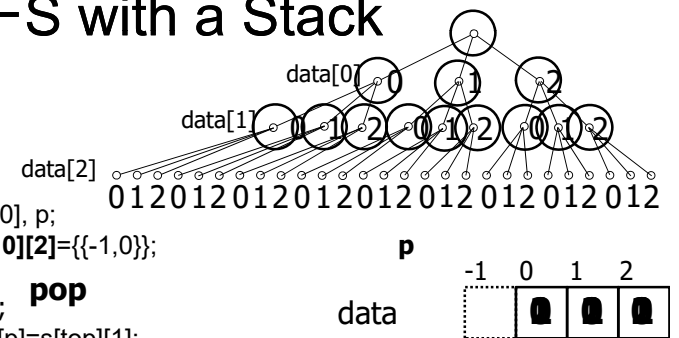
Recursive DFS implementation

```
07 void countUp(int n, int data[], int m, int pivot) {
08     if (pivot>=m)
09         print(data, m);
10     else
11         for (data[pivot]=0; data[pivot]<n; data[pivot]++)
12             countUp(n, data, m, pivot+1);
13 }
```

11

DFS with a Stack

```
01 #include <stdio.h>
02 int main() {
03     int n=3, i, data[10], p;
04     int top=1, s[10*10][2]={{-1,0}};
05     while (top>0) {
06         pop
07         p=s[--top][0];
08         if (p>=0) data[p]=s[top][1];
09         if (p==n-1)
10             for (i=0; i<n; i++)
11                 printf("%d%c", data[i], i<n-1?' ':'\n');
12         else
13             for (p++; p<n; p++)
14                 s[top][0]=p, s[top][1]=i, top++;
15     }
16     return 0;
17 }
```



2, 0
2, 0
2, 0
2, 0
0, 0
0, 2
-1, 0

12

Permutation

- If you view the index[] as an n-digit n-ary number (the digits set is {1, 2, ..., n}), the previous program generates all possible increasing numbers (from 111...1 to nnn...n) and eliminates all invalid numbers with repeating digits on the way of counting up.
- Problem: slow as n gets to 10 or larger. n^n vs. $n!$
- Next, you will see another implementation. It tries to generate only valid increasing numbers with the next() function. Later, we will have some more algorithms to generate permutations.

13

```

01 #include <stdio.h>
02
03 void main()
04 {
05     int size, perm[12] = {0}, current=0, solCount=0, i;
06
07     printf("Please input number of elements: ");
08     scanf("%d", &size);
09
10     while (current >= 0)
11     {
12         current += next(size, current, perm);
13         if (current == size)
14         {
15             solCount++;
16             printf("%4d: ", solCount);
17             for (i=0; i<size; i++)
18                 printf("%d ", perm[i]);
19             printf("\n");
20             current = size-1;
21         }
22     }
23     printf("Total %d permutations\n", solCount);
24 }

```

perm

1	2			
---	---	--	--	--

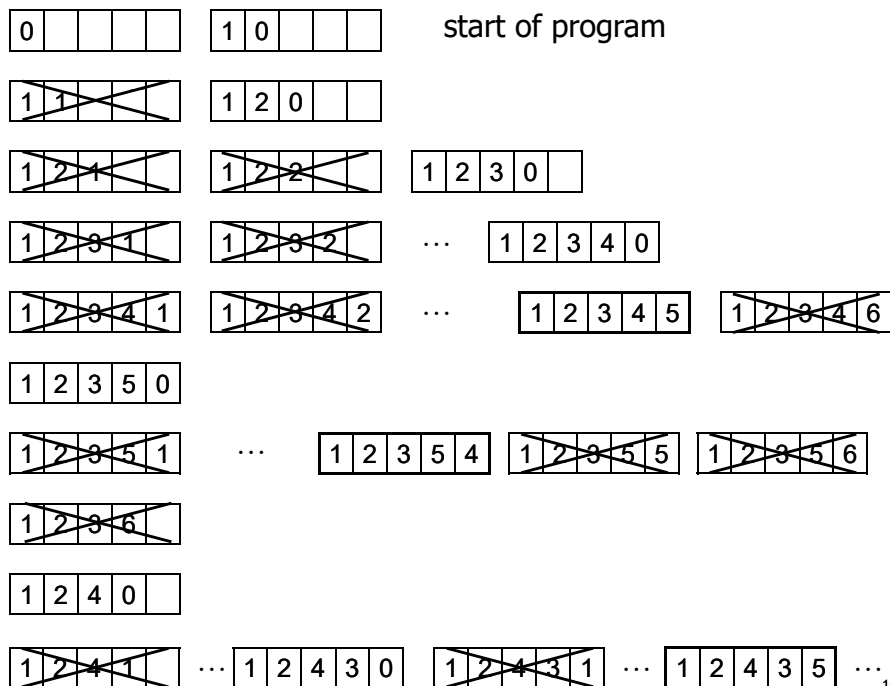
current

```

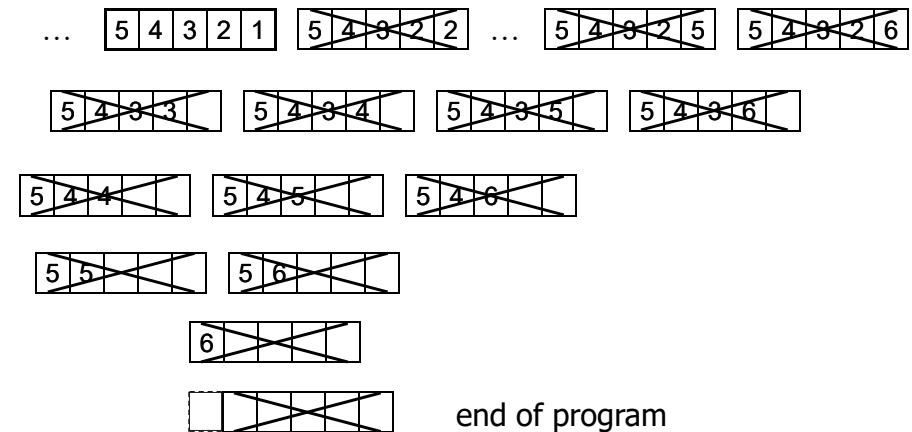
26 int next(int size, int pivot, int perm[])
27 {
28     int i, collision;
29
30     while (perm[pivot]++ < size)
31     {
32         collision = 0;
33         for (i=0; i<pivot; i++)
34             if (perm[pivot] == perm[i])
35                 collision = 1;
36         if (collision) break;
37     }
38     if (!collision) return 1;
39 }
40 perm[pivot] = 0;
41 return -1;
42 }
43 }

```

14



15



16

Permutation from Swapping

• Recursive

- 1 + permutations of {2, 3, 4}
- 2 + permutations of {1, 3, 4}
- 3 + permutations of {2, 1, 4}
- 4 + permutations of {2, 3, 1}

```
for (i=0; i<n; i++) a[i] = i+1;
permutation(a, 0, n-1);
```

```
void permutation(int perm[], int start, int end) {
    if (start == end) { printPerm(perm, end+1); return; }
    for (int i=start; i<=end; i++) {
        swap(&perm[start], &perm[i]);
        permutation(perm, start+1, end);
        swap(&perm[start], &perm[i]);
    }
}
```

1	2	3	4
1	2	4	3
1	3	2	4
1	3	4	2
1	4	3	2
1	4	2	3
2	1	3	4
2	1	4	3
...	...	...	...
3	2	1	4
3	2	4	1
...	...	...	...
4	2	3	1
4	2	1	3
...	...	...	...
4	1	2	3

from Swapping (cont'd) ZJ e446

- The output of previous program is **not** in **lexicographic order**

```
1: 1 2 3 4
2: 1 2 4 3
3: 1 3 2 4
4: 1 3 4 2
5: 1 4 3 2
6: 1 4 2 3
7: 2 1 3 4
8: 2 1 4 3
9: 2 3 1 4
10: 2 3 4 1
11: 2 4 3 1
12: 2 4 1 3
5: 1 4 2 3
6: 1 4 3 2
```

```
13: 3 1 2 4
14: 3 1 4 2
15: 3 2 1 4
16: 3 2 4 1
17: 3 4 1 2
18: 3 4 2 1
19: 4 1 2 3
20: 4 1 3 2
21: 4 2 1 3
22: 4 2 3 1
23: 4 3 1 2
24: 4 3 2 1
```

```
for (int i=start; i<=end; i++) {
    swap(&perm[start], &perm[i]);
    permutation(perm, start+1, end);
    swap(&perm[start], &perm[i]);
}
```

如果排列的數字允許重複?
ITSA33 problem #2

start → i → end
perm[] 3 2 1 4

```
void pRotate(int a, int b) {
    int tmp=a, d=b>a?-1:1;
    while (a!=b) *b = *(b+d), b+=d;
    *a = tmp;
}
```

18

Permutation from Rotation

0 1 2 3	2 0 1 3	1 2 0 3	1 0 2 3	2 1 0 3	0 2 1 3
3 0 1 2	3 2 0 1	3 1 2 0	3 1 0 2	3 2 1 0	3 0 2 1
2 3 0 1	1 3 2 0	0 3 1 2	2 3 1 0	0 3 2 1	1 3 0 2
1 2 3 0	0 1 3 2	2 0 3 1	0 2 3 1	1 0 3 2	2 1 3 0
0 1 2 3	2 0 1 3	1 2 0 3	1 0 2 3	2 1 0 3	0 2 1 3
		0 1 2 3			1 0 2 3
					0 1 2 3

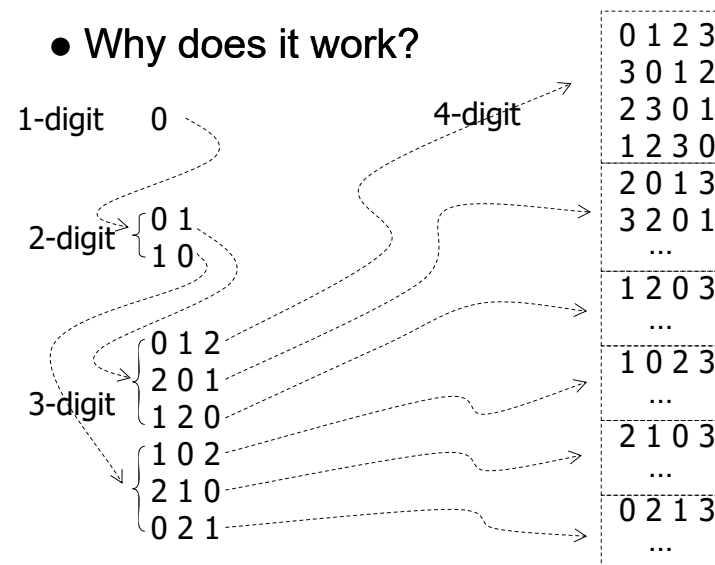
```
int rotate_n_check(int size, char data[]) {
    int i;
    for (i=size; i>=2; i--) {
        rotate(i, data);
        if (data[i-1] != i-1) return 1;
    }
    return 0;
}
```

```
for (i=0; i<num; i++) digit[i] = i;
do
    printPermutation(num, digit);
while (rotate_n_check(num, digit));
```

19

Rotation (cont')

- Why does it work?



20

Recursive Implementation

```

void permuteFromRotation(char pdata[], int len) {
    char cdata[20];
    int i, j, tmp;
    if (len >= n) {
        pdata[n] = 0; printf("%s\n", pdata);
        return;
    }
    for (i=0; i<len; i++) cdata[i] = pdata[i];
    cdata[len] = items[len];
    for (i=0; i<=len; i++) {
        permuteFromRotation(cdata, len+1);
        for (tmp=cdata[len], j=len; j>0; j--)
            cdata[j] = cdata[j-1];
        cdata[0] = tmp;
    }
}

result[0] = items[0];
permuteFromRotation(result, 1);
    
```

Diagram illustrating the recursive process for generating permutations. The array `items` contains the sequence '0', '1', '2', '3'. The function `permuteFromRotation` is called with `len` values from 1 to 4. The diagram shows the state of `pdata` and `cdata` at each step, with arrows indicating the rotation of elements in `cdata` to generate the next permutation in `pdata`.

21

Recursive Implementation

```

char items[] = "0123"; int n=4;
void permuteFromRotation(char pdata[], int len) {
    int i, j, tmp;
    if (len >= n) {
        pdata[n] = 0; printf("%s\n", pdata);
        return;
    }
    pdata[len] = items[len];
    for (i=0; i<=len; i++) {
        permuteFromRotation(pdata, len+1);
        for (tmp=pdata[len], j=len; j>0; j--)
            pdata[j] = pdata[j-1];
        pdata[0] = tmp;
    }
}

result[0] = items[0];
permuteFromRotation(result, 1);
    
```

Diagram illustrating the recursive process for generating permutations. The array `items` contains the sequence '0', '1', '2', '3'. The function `permuteFromRotation` is called with `len` values from 1 to 4. The diagram shows the state of `pdata` and `cdata` at each step, with arrows indicating the rotation of elements in `pdata` to generate the next permutation in `pdata`.

22

Permutation by Exchanging Neighbors

```

perm pos perm pos
3 2 1 0 0 0 0 3 2 0 1 0 0 1
2 3 1 0 1 0 0 2 3 0 1 1 0 1
2 1 3 0 2 0 0 2 0 3 1 2 0 1
2 1 0 3 3 0 0 2 0 1 3 3 0 1
3 1 2 0 0 1 0 3 0 2 1 0 1 1
1 3 2 0 1 1 0 0 3 2 1 1 1 1
1 2 3 0 2 1 0 0 2 3 1 2 1 1
1 2 0 3 3 1 0 0 2 1 3 3 1 1
3 1 0 2 0 2 0 3 0 1 2 0 2 1
1 3 0 2 1 2 0 0 3 1 2 1 2 1
1 0 3 2 2 2 0 0 1 3 2 2 2 1
1 0 2 3 3 2 0 0 1 2 3 3 2 1

for (i=0; i<num; i++)
    perm[i] = num-1-i, position[i] = 0;
do
    printPermutation(num, perm);
while (shift_n_check(num, perm, position))

void shiftRight(int size, char perm[], char *pos){
    char tmp = perm[*pos];
    if (*pos < size-1) {
        perm[*pos] = perm[*pos+1];
        perm[++(*pos)] = tmp;
    }
    else {
        for (int i=size-2; i>=0; i--)
            perm[i+1] = perm[i];
        perm[*pos=0] = tmp;
    }
}

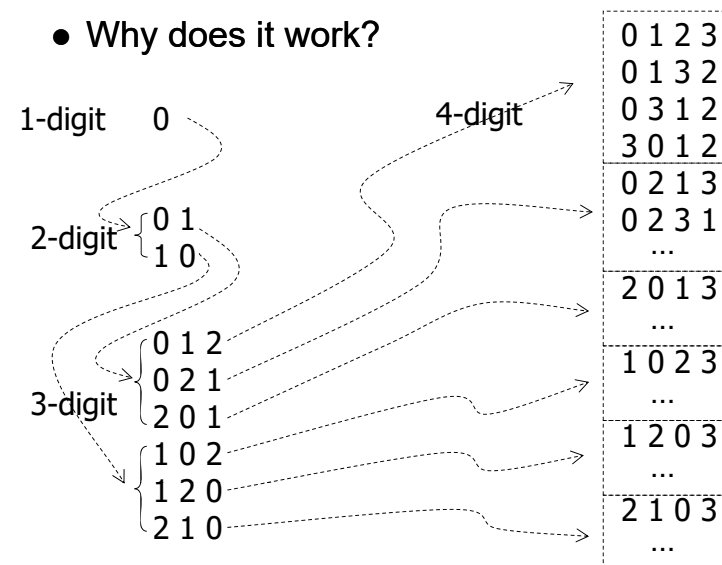
int shift_n_check(int size, char perm[], char pos[]){
    for (int i=size; i>=2; i--) {
        shiftRight(i, &perm[size-i], &pos[size-i]);
        if (pos[size-i] > 0) return 1;
    }
    return 0;
}
    
```

Diagram illustrating the permutation process by exchanging neighbors. The diagram shows the state of the permutation array `perm` and the position array `pos` at each step, with arrows indicating the shift of elements in `perm` to generate the next permutation in `perm`.

23

Exchanging Neighbors (cont'd)

- Why does it work?



24

Recursive Implementation

```
void permuteFromExchanging(char pdata[], int len) {
    char cdata[10];
    int i, j, tmp;
    if (len >= n) {
        pdata[n] = 0; printf("%s\n", pdata);
        return;
    }
    for (i=0; i<len; i++) cdata[i] = pdata[i];
    cdata[len] = items[len];
    for (i=len; i>=0; i--) {
        permuteFromExchanging(cdata, len+1);
        if (i>0)
            tmp=cdata[i], cdata[i]=cdata[i-1], cdata[i-1] = tmp;
    }
    result[0] = items[0];
    permuteFromExchanging(result, 1);
}
```

len 2

pdata[]

cdata[]

items[]

n 4

25

Recursive Impl. w/o Extra Array

```
void permuteFromExchanging(char data[], int len) {
    int i, j, tmp;
    if (len >= n) {
        data[n] = 0; printf("%s\n", data);
        return;
    }
    data[len] = items[len];
    for (i=len; i>=0; i--) {
        permuteFromExchanging(data, len+1);
        if (i>0)
            tmp=data[i], data[i]=data[i-1], data[i-1] = tmp;
    }
    for (i=0; i<len; i++) data[i] = data[i+1];
    result[0] = items[0];
    permuteFromExchanging(result, 1);
}
```

len 2

data[]

items[]

n 4

26

Generate Disjoint Partitions

- A partition of a set X is a set of non-empty subsets of X such that every element x in X is in exactly one of these subsets.

- e.g. The set X={1, 2, 3} has five partitions:

- { {1}, {2}, {3} }, sometimes written 1 | 2 | 3.
- { {1, 2}, {3} }, or 12 | 3.
- { {1, 3}, {2} }, or 13 | 2.
- { {1}, {2, 3} }, or 1 | 23.
- { {1, 2, 3} }, or 123

- The total number of partitions is the Bell numbers B_n

$$B_{n+1} = \sum_{k=0}^n C_k^n B_k$$

e.g. $B_0 = 1, B_1 = 1, B_2 = 2, B_3 = 5, B_4 = 15, B_5 = 52, B_6 = 203, \dots$

$B_{20} = 51724158235372 \approx 5 \times 10^{13}, B_{30} = 846749014511809332450147 \approx 8 \times 10^{23},$

$B_{40} = 157450588391204931289324344702531067 \approx 2 \times 10^{35},$

$B_{50} = 185724268771078270438257767181908917499221852770 \approx 2 \times 10^{47}.$

27

Generate Partitions (cont'd)

- A simple algorithm through counting up with dynamically adjusted ceiling, partition / max e.g., X is a 4-element set $\{a_0, a_1, a_2, a_3\}$

- 0,0,0,0 / 0,1,1,1 $\Rightarrow \{a_0, a_1, a_2, a_3\}$
- 0,0,0,1 / 0,1,1,1 $\Rightarrow \{a_0, a_1, a_2\}, \{a_3\}$
- 0,0,1,0 / 0,1,1,2 $\Rightarrow \{a_0, a_1, a_3\}, \{a_2\}$
- 0,0,1,1 / 0,1,1,2 $\Rightarrow \{a_0, a_1\}, \{a_2, a_3\}$
- 0,0,1,2 / 0,1,1,2 $\Rightarrow \{a_0, a_1\}, \{a_2\}, \{a_3\}$
- 0,1,0,0 / 0,1,2,2 $\Rightarrow \{a_0, a_2, a_3\}, \{a_1\}$
- 0,1,0,1 / 0,1,2,2 $\Rightarrow \{a_0, a_2\}, \{a_1, a_3\}$
- 0,1,0,2 / 0,1,2,2 $\Rightarrow \{a_0, a_2\}, \{a_1\}, \{a_3\}$
- 0,1,1,0 / 0,1,2,2 $\Rightarrow \{a_0, a_3\}, \{a_1, a_2\}$
- 0,1,1,1 / 0,1,2,2 $\Rightarrow \{a_0\}, \{a_1, a_2, a_3\}$
- 0,1,1,2 / 0,1,2,2 $\Rightarrow \{a_0\}, \{a_1, a_2\}, \{a_3\}$
- 0,1,2,0 / 0,1,2,3 $\Rightarrow \{a_0, a_3\}, \{a_1\}, \{a_2\}$
- 0,1,2,1 / 0,1,2,3 $\Rightarrow \{a_0\}, \{a_1, a_3\}, \{a_2\}$
- 0,1,2,2 / 0,1,2,3 $\Rightarrow \{a_0\}, \{a_1\}, \{a_2, a_3\}$
- 0,1,2,3 / 0,1,2,3 $\Rightarrow \{a_0\}, \{a_1\}, \{a_2\}, \{a_3\}$

Implementation with 2 arrays

0,0,1,2 label the partitions for each set elements a_i

partition

0	0	1	2
---	---	---	---

max

0	1	1	2
---	---	---	---

each $\max[i] \geq$ all prior partitions

when a non-tailing digit reaches its max

partition

0	1	2	0
---	---	---	---

max

0	1	2	3
---	---	---	---

28

Implementation

```

01 int next(int size, int pivot, int partition[], int max[]) {
02     int i, j, maxVal;
03     while (pivot>0 && (partition[pivot] >= max[pivot]))
04         pivot--;
05     if (pivot > 0) {
06         partition[pivot]++;
07         for (i=pivot+1; i<size; i++) {
08             partition[i] = 0;
09             maxVal = 0;
10             for (j=1; j<i; j++)
11                 if (partition[j]>maxVal)
12                     maxVal = partition[j];
13             max[i] = maxVal+1;
14         }
15         return size;
16     }
17     return pivot;
18 }

```

pivot ↘

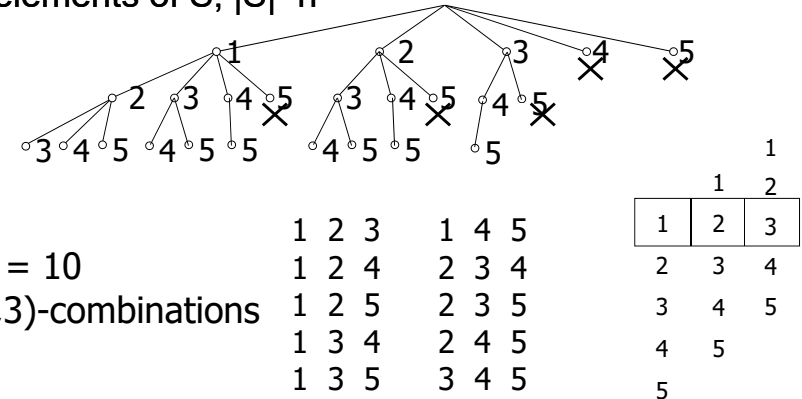
	0	1	2	3
partition	0	1	1	2
max	0	1	2	2

	0	1	2	3
partition	0	1	2	0
max	0	1	2	3

29

Generate (n,k)-Combinations

- (n,k)-combination of a set S is a subset of k distinct elements of S, |S|=n



$C_3^5 = 10$
(5,3)-combinations

- Extend the counting-up program and avoid those non-increasing sequences

30

Efficient Enumerating w/o Invalid()

```

int next(int comb[], int n, int k) {
    int i, j;
    for (i=k-1; i>=0; i--)
        if (comb[i]<n-(k-1-i)) {
            comb[i]++;
            for (j=i+1; j<k; j++)
                comb[j] = comb[j-1]+1;
            return 1;
        }
    return 0;
}

```

comb

1	2	3
1	2	4
1	2	5
1	3	4
1	3	5
1	4	5
2	3	4
2	3	5
2	4	5
3	4	5

```

int comb[10], n=5, k=3;
for (i=0; i<k; i++)
    comb[i] = i+1;
do
    printArray(comb, k);
while (next(comb, n, k));

```

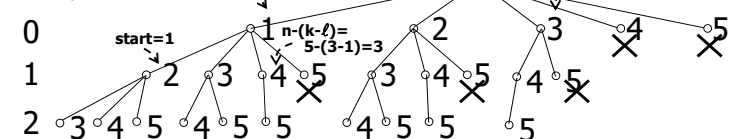
Output symbols

'Z' 'A' 'K' 'D' 'S'

31

Recursive (n,k)-Combinations

level- ℓ $C(5,3)$ start=0



- select level- ℓ elements from seq[start]~seq[n-(k- ℓ)]
- put the result to result[ℓ]~result[k-1]
- comb(n,k,seq,0,0,result); // DFS backtracking

```

void comb(int n, int k, const int seq[], int start, int l, int result[]) {
    if (l==k) print(result, k); // result[0]~result[k-1]
    else
        for (int i=start; i<=n-(k-l); i++)

```

e.g. brute force UVa10326
UVa11659 w/ binary search

```

        result[l] = seq[i];
        comb(n, k, seq, i+1, l+1, result);
    }
}

```

level result 1 2 3 seq 1 2 3 4 5

32

Generate Power Set

- The power set of a set is the collection of all subsets of S, including the empty set and S itself.
 - e.g. $S = \{1, 2, 3\}$,
 $2^S = \{\{\}, \{1\}, \{2\}, \{3\}, \{1,2\}, \{1,3\}, \{2,3\}, \{1,2,3\}\}$
- We can extend the program for generating (n,k)-combinations to generate the whole power set.
- int n=3, seq[]={1,2,3}, result[3], k;
 for (k=0; k<=n; k++)
 comb(n, k, seq, 0, 0, result);

33

正整數 N 的加法拆解 (整數劃分)

- 每個正整數 N 都可以寫成比它小的數字和, 請把所有不重複整數加總等於 N 的方法寫出來

Sample Input 6	Sample Input 10	Sample Input 15	
Sample Output 1 2 3 1 5 2 4 6	Sample Output 1 2 3 4 1 2 7 1 3 6 1 4 5 2 3 5 3 7 4 6 10	Sample Output 1 2 3 4 5 1 2 3 9 1 2 4 8 1 2 5 7 1 2 12 1 3 4 7 1 3 5 6 1 3 11 1 4 10 1 5 9 1 6 8 1 14	2 3 4 6 2 3 10 2 4 9 2 5 8 2 6 7 2 13 3 4 8 3 5 7 3 12 4 5 6 4 11 5 10 6 9 7 8 15

- 輸出請依照範例以一般化字典序 (lexicographic order) 排列

34

Sudoku

- Sudoku:** In these three examples, 81 cells are divided into 9 blocks each with 9 cells (3-by-3). A player is required to fill in the blank cells such that integers in each row, each column, and each block are permutations of $\{1,2,3,\dots,9\}$, i.e. no duplication of numbers in each row, column, or block.

7	8	9	2		5
6			5		8
				1	
2		8		5	1
	5	7	1	6	
9	1		6		3
	7				
	5	9			6
8			1	5	3

number of lines
row, column, value
row, column, value
...

Initial configuration

30
0, 0, 7
0, 1, 8
0, 2, 9
0, 4, 2
0, 8, 5
1, 0, 6
1, 5, 5
1, 7, 8
...

	6	1	3	4	
	3		0		
5	4	7		1	2
		2			4
	5	3	9	7	
7			4		
	3	7		5	4
		8		7	
	1	4	7	8	

3	8			6	
6				4	
9	8	5		1	
2		9			
4	5			9	2
			7		6
2		1	3	6	
	9				3
	1			5	4

35

Sudoku (cont'd)

- Extension of counting

2	7	6	1		3	4		
1	9	3			8			
5	4							
	5							
7								
	3							

...

✕	✕	...	1	✕
2	✕		✕	✕
✕	✕		✕	✕
✕	✕		✕	✕
✕	✕		✕	✕
✕	✕		✕	✕
✕	7		✕	✕
8	8		✕	✕
9	9		9	9

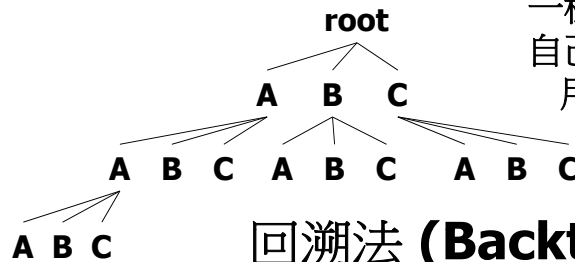
		6	1		3	4		
		3			8			
5	4	7			1	2		
		2					4	
	5	3	9		7			
7			4					
	3	7		5		4	2	
		8			7			
	1	4	7	8				

more constraints
on the set of
values to be filled
in each cell

36

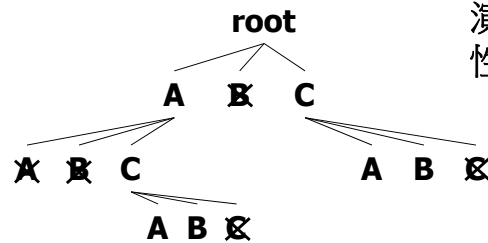
遞迴函式 (Recursive Function)

一種程式實作方式
自己呼叫自己，常常
用來窮舉所有的
可能性, DFS



回溯法 (Backtracking)

演算法: 窮舉所有的可能性, 但是一發現不行趕快
回頭嘗試下一個可能的
路徑, 常常用遞迴
函式來實作, 也可以
用迴圈實作



12-1 尋找整除數字 12-2 速讀

37

Backtracking with Iteration

constraints: DFS → backtracking

```
01 void sudoku(int board[], int seq[], int nseq, int p) {
02     while (p >= 0)
03         if (p >= nseq) print_solution, p--;
04         else if (board[seq[p]] >= 9)
05             board[seq[p]] = 0;
06         else {
07             board[seq[p]]++;
08             if (isValid(board, seq[p])) p++;
09         }
10 }
```

6	1	3	4		
	3		8		
5	4	7		1	2
		2			4
5	3	9	7		
7			5		
3	7		4	4	2
		8		7	
	1	4	7	8	

30
0, 0, 7
0, 1, 8
0, 2, 9
0, 4, 2
0, 8, 5
1, 0, 6
1, 5, 5

board seq
0 ← 0
0 ← 1
6 ← 4
1 ← 7
0 ← 8
3 ← 9
4 ← 10
0 ← ...
0 ← 79
... ← 80
8 ← ...
0 ← ...
0 ← ...
nseq = 51

```
01 int isValid(int data[], int p) {
02     int i, j, r = p/9, c = p%9;
03     for (i = 0; i < 9; i++) {
04         if (i != c && data[p] == data[r*9+i]) return 0;
05         if (i != r && data[p] == data[i*9+c]) return 0;
06     }
07     for (i = r/3*3; i < r/3*3+3; i++)
08         for (j = c/3*3; j < c/3*3+3; j++)
09             if ((i != r || j != c) && data[p] == data[i*9+j]) return 0;
10     return 1;
11 }
12 }
```

```
01 int main() {
02     int board[81] = {};
03     int seq[81], nseq = 0;
04     read constraints to board[81]
05     prepare seq[81] and nseq
06     sudoku(board, seq, nseq, 0);
07     return 0;
08 }
```

38

Backtracking using Recursion

without any constraint, raw DFS to generate 9^{81} possibilities

```
01 int main() {
02     int data[81]; // 9x9
03     sudoku(data, 0);
04     return 0;
05 }

01 void print(int data[]) {
02     int i, j;
03     for (i = 0; i < 9; i++)
04         for (j = 0; j < 9; j++)
05             printf("%2d%s", data[i*9+j], j == 9 ? "\n" : "");
06 }
```

```
01 void sudoku(int data[], int p) {
02     if (p >= 81)
03         print(data);
04     else
05         for (data[p] = 1; data[p] <= 9; data[p]++)
06             sudoku(data, p+1);
07 }
```

39

Recursion (cont'd)

constraints: DFS → backtracking

```
01 void sudoku(int board[], int seq[], int nseq, int p) {
02     if (p >= nseq)
03         print(board);
04     else
05         for (board[seq[p]] = 1; board[seq[p]] <= 9; board[seq[p]]++)
06             if (isValid(board, seq[p]))
07                 sudoku(board, seq, nseq, p+1);
08 }
```

```
01 int isValid(int data[], int p) {
02     int i, j, r = p/9, c = p%9;
03     for (i = 0; i < 9; i++) {
04         if (i != c && data[p] == data[r*9+i]) return 0;
05         if (i != r && data[p] == data[i*9+c]) return 0;
06     }
07     for (i = r/3*3; i < r/3*3+3; i++)
08         for (j = c/3*3; j < c/3*3+3; j++)
09             if ((i != r || j != c) && data[p] == data[i*9+j])
10                 return 0;
11     return 1;
12 }
```

6	1	3	4		
	3		8		
5	4	7		1	2
		2			4
5	3	9	7		
7			5		
3	7		4	4	2
		8		7	
	1	4	7	8	

30
0, 0, 7
0, 1, 8
0, 2, 9
0, 4, 2
0, 8, 5
1, 0, 6
1, 5, 5

board seq
0 ← 0
0 ← 1
6 ← 4
1 ← 7
0 ← 8
3 ← 9
4 ← 10
0 ← ...
0 ← 79
... ← 80
8 ← ...
0 ← ...
0 ← ...
nseq = 51

```
01 int main() {
02     int board[81] = {};
03     int seq[81], nseq = 0;
04     read constraints to board[81]
05     prepare seq[81] and nseq
06     sudoku(board, seq, nseq, 0);
07     return 0;
08 }
```

40

5	3	1	2	7	6	8	9	4
6	2	4	1	9	5	2		
	9	8					6	
8				6				3
4			8		3			1
7				2				6
	6					2	8	
			4	1	9			5
				8			7	9

41

Determinant of an n-by-n Matrix

- Leibniz formula

$$\mathbf{A} = \begin{pmatrix} a_{1,1} & a_{1,2} & a_{1,3} \\ a_{2,1} & a_{2,2} & a_{2,3} \\ a_{3,1} & a_{3,2} & a_{3,3} \end{pmatrix}$$

- Determinant

$$\det(\mathbf{A}) = a_{1,1}a_{2,2}a_{3,3} + a_{1,2}a_{2,3}a_{3,1} + a_{1,3}a_{2,1}a_{3,2} - a_{1,3}a_{2,2}a_{3,1} - a_{1,2}a_{2,1}a_{3,3} - a_{1,1}a_{2,3}a_{3,2}$$

- Permutation?

$$\det(\mathbf{A}) = a_{1,1}a_{2,2}a_{3,3} + a_{1,2}a_{2,3}a_{3,1} + a_{1,3}a_{2,1}a_{3,2} - a_{1,3}a_{2,2}a_{3,1} - a_{1,2}a_{2,1}a_{3,3} - a_{1,1}a_{2,3}a_{3,2}$$

- $\text{sign}(\sigma) = \begin{cases} 1 & \text{if } \sigma \text{ has even number of misplaced pairs} \\ -1 & \text{if ... odd ...} \end{cases}$

e.g. $\text{sign}(123)=1$, $\text{sign}(231)=1$, $\text{sign}(312)=1$
 $\text{sign}(321)=-1$, $\text{sign}(213)=-1$, $\text{sign}(132)=-1$

42

Applications

1. 給定 n 個各種長度的木棒, 是否可排出正方形? 是否可排出指定長寬比例的矩形?
2. 有 m 條星狀連接的路徑, 在各路徑途中指定位置有不同金額的獎勵, 如果油料有限制, 最多只能夠走 n 公里, 請找到由中心點出發能夠取回最大獎勵的方式
3. n 個小偷竊得 m 個不同價值的物品, 怎樣才能比較公平地分贓而不至於內鬨呢?
4. n 組數字, 由每一組數字中各取一個, 總和最大的取法?
5. 尋找 $n \times n$ 魔方陣, 洛書
6. Tetromino & pentomino
7. 輸入整數 N , $2 \leq N \leq 9$, 印出一 N 位數字 X , 其 N 個數字都不重複且 $\text{mod } 10^i$ 皆能被 N 整除, 例如 $N = 8$, $X = 21579648$ 、 $X \text{ mod } 10^7 = 1579648$ 、 579648 、 79648 、 9648 、 648 、 48 、 8 皆可以被 8 整除且各位數字不重複出現。

43